

# A Python Implementation of Chebyshev Functions

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# A Quick Note

**To the applied mathematicians:** I know Nick Trefethen gave several talks on Chebfun this year. Don't worry! I hope to discuss some things he didn't talk about:

- why his methods works so well,
- how to actually compute Chebfun,
- how to actually compute integrals, derivatives, global roots.

# Two Key Theorems

## Weierstrass Approximation Theorem on $\mathbb{R}$ (1885)

Suppose  $f \in C^0[a, b]$ . Then  $\forall \epsilon > 0, \exists p$  polynomial such that  $\|f - p\|_\infty < \epsilon$  on  $[a, b]$ .

(Interesting corollary:  $\mathbb{R}[x]$  is dense in  $C^0[a, b]$ .  $\mathbb{Q}[x]$  is dense in  $\mathbb{R}[x]$ . Therefore  $|C^0[-1, 1]| = |\mathbb{R}|$ .)

## Restriction to Degree $n$ Polynomials: $\mathcal{P}_n$

Given  $f \in C^0[a, b]$  find  $p^* \in \mathcal{P}_n$  such that

$$\|f - p^*\|_\infty \leq \|f - p\|_\infty \quad \text{for all } p \in \mathcal{P}_n$$

**Fact:**  $p^*$  exists, is unique, and goes by the name “best”, “uniform”, “Chebyshev”, or “minimax” approximation to  $f$ .

# Objective

## Main Goal

Given  $f : [-1, 1] \rightarrow \mathbb{R}$ , find a polynomial of lowest degree that approximates  $f$  to within machine epsilon. That is, find

$$p_{\text{goal}} = \min_{p, \deg(p)} \|f - p\|_{\infty} < \epsilon_{\text{mach}}$$

Trefethen's method: look at polynomials with a particular structure that make derivatives, integrals, and roots easy to compute.

Use these tools:

- Structure #1: Lagrange interpolants
- Structure #2: Chebyshev polynomial expansions

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## Adding Some Structure

Trefethen restricts to the space of Lagrange interpolating polynomials. Here we will show close they get to the minimax polynomial approximation.

### Lagrange Interpolating Approximations

Let  $L_n(f)$  be a Lagrange interpolant of  $f$  on  $n$  nodes. Then,

$$\|f - L_n(f)\|_\infty \leq (1 + \|L_n\|_\infty) \|f - p^*\|_\infty$$

where  $p^*$  is the minimax polynomial approximation of degree  $n$ .

$\|L_n\|$  depends on the choice of interpolating points:

- Uniform Distribution:  $\|L_n\|_\infty = O\left(\frac{2^{n+1}}{n \log n}\right)$
- Chebyshev Distribution:  $\|L_n\|_\infty = O(\log(n+1))$



# More Reasons to Use Lagrange Interpolation

Lagrange Interpolants over the Chebyshev points have good convergence properties:

## Bounded Variation

If  $\partial^k f : [-1, 1] \rightarrow \mathbb{R}$  has bounded variation for some  $k \geq 1$  then

$$\|f - L_n(f)\| = O(n^{-k})$$

## Analyticity

If  $f$  is analytic in a neighborhood of  $[-1, 1]$  then

$$\|f - L_n(f)\| = O(C^n)$$

for some  $C < 1$ .

## Conclusions About Accuracy

### Theorem

Chebyshev interpolants are *near-best* or *spectrally accurate*.

*Spectral accuracy* has to do with how well an approximating function converges in a spectral domain. (i.e. Fourier) The theory of Sobolev spaces gives a rich and precise description of why this works. I would love to discuss this, but I won't.

**The point:** Lagrange interpolants over the Chebyshev points (Chebyshev interpolants) are good polynomial approximations.

- For example: they defeat the Gibbs phenomenon.

# Practical Uses of Lagrange Interpolation

## Barycentric Formula

Let  $(x_j, f_j)$  be a collection of  $N + 1$  sample points of  $f$ . Then the Lagrange interpolant over these points can be written as

$$p(x) = \frac{\sum_{j=0}^N \frac{w_j}{x - x_j} f_j}{\sum_{j=0}^N \frac{w_j}{x - x_j}}$$

with  $w_j = (-1)^j$ . (Divide by two for  $j = 0, N$ .)

Benefits:

- fast ( $O(N)$  evaluation operation count),
- numerically stable,
- $O(N)$  “updating” method for grid refinements.

## Why Restriction to $[-1, 1]$ ?

Summary so far: we've seen the benefits of using Lagrange interpolation over the Chebyshev points and how the Barycentric formula allows  $O(N)$  evaluation as well as  $O(N)$  method of adding Chebyshev points.

Trefethen exploits another structure: consider only  $f : [-1, 1] \rightarrow \mathbb{R}$

### Chebyshev Polynomials

The Chebyshev polynomials  $T_n : [-1, 1] \rightarrow \mathbb{R}$  are defined:

$$T_0(x) = 1$$

$$T_1(x) = x$$

$$T_{n+1}(x) = 2xT_n(x) - T_{n-1}(x)$$

# Chebyshev Polynomial Expansions

## Theorem (Trefethen)

If  $f : [-1, 1] \rightarrow \mathbb{R}$  is Lipschitz continuous then the Chebyshev expansion

$$g(x) = \sum_{n=0}^{\infty} a_n T_n(x)$$

converges absolutely and uniformly to  $f$ . Additionally, the Chebyshev coefficients are given by

$$a_n = \frac{2}{\pi} \int_{-1}^1 \frac{f(x) T_n(x)}{\sqrt{1-x^2}} dx$$

with the special case that for  $n = 0$  the constant changes to  $1/\pi$ .

# Relating Chebyshev Polynomials to Chebyshev Interpolants

## Chebyshev Polynomial Representations

For  $x \in [-1, 1]$  define  $\theta = \arccos(x) \in [0, 2\pi]$  and  $z = e^{i\theta}$ . Then:

$$T_n(x) = \cos n\theta = \frac{1}{2} (z^n + z^{-n})$$

For Chebyshev interpolants, this is a finite sum:

## Expansion Representation

Let  $p$  be a Chebyshev interpolant determined by  $N + 1$  grid points  $(x_i, f_i)$  over the Chebyshev points  $x_i = \cos(\pi i/N)$ . Then

$$p(x) = \sum_{n=0}^N a_n T_n(x) = \sum_{n=0}^N a_n \cos(n\theta) = \frac{1}{2} \sum_{n=0}^N a_n (z^n + z^{-n})$$

# Definition of Chebfun

## Definition

A **chebfun** is a minimal degree Barycentric Lagrange interpolant of a function  $f : [-1, 1] \rightarrow \mathbb{R}$  over the Chebyshev points that uses its Chebyshev polynomial expansion for fast computations.

- “Spectrally optimal approximate”
- Next section: Given  $f \in Lip[-1, 1]$ , compute its Chebfun,  $p$ .

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# Why Python?

- Python is a popular, powerful, and easy to use programming language. Moreover, it's open-source and free!
  - Numpy/Scipy
  - Sage
  - Clawpack
  - Brian
- Support open-source! (<http://www.opensource.org>)



# Generating Chebyshev Functions

- Init: begin with  $N = 4$  Chebyshev interpolating points,  $\{x_i\}$  and compute  $f_i = f(x_i)$ .
- Loop: Compute the discrete cosine transform of the  $\{f_i\}$ ,  $\{\hat{f}_i\}$ , and divide each term by the number of interpolating points,  $N$ . Set  $a_i = \hat{f}_i/N$ .

- **If:**

$$|a_N|, |a_{N-1}| < 2 * \epsilon_{\text{mach}} * \max_i |a_i|$$

then truncate the sequence  $\{a_i\}$  down to the term with largest index  $M$  exceeding this bound. The remaining  $M$  terms gives you the “optimal Chebyshev interpolant size.

- **Else:** Loop with  $2N$ .

# *Demo*

# Integration

## Chebyshev Polynomial Integration

$$\int_{-1}^1 T_n(x) dx = \begin{cases} 0 & \text{if } n \text{ odd} \\ \frac{2}{1-n^2} & \text{if } n \text{ even} \end{cases}$$

Integral of a chebfun on  $[-1, 1]$ :

## Clenshaw-Curtis Quadrature

Given a chebfun  $p(x) = \sum_{n=0}^N a_n T_n(x)$ ,

$$\int_{-1}^1 p(x) dx = \sum_{n \text{ even}}^N \frac{2}{1-n^2}$$

# Differentiation

## Chebyshev Polynomial Derivative

If  $p(x) = \sum_{n=0}^N a_n T_n(x)$  then  $p'(x) = \sum_{n=0}^{N-1} b_n T_n(x)$  with

$$b_{n-1} = b_{n+1} + 2na_n, \quad b_N = b_{N+1} = 0, \quad b_0 = b_2/2 + a_1.$$

- Backsolve to find  $b_n$  and inverse cosine transform to find  $(x_i, f_i)$  pairs.
- $O(N)$  computation of function derivative.

# Root Finding

Recall that Chebyshev coefficients  $\{a_n\}$  are the coefficients of the Laurent series

$$q(z) = \sum_{n=0}^N a_n (z^n + z^{-n}).$$

- The roots,  $z_i$  of  $q$ , are the roots of  $z^N q(z)$ ; a polynomial of degree  $2N$  in  $z$ .
- Find the roots of  $z^N q(z)$  using your favorite polynomial root-finding algorithm. (Much easier than for a general function.)
- Roots come in pairs  $\{z_n, z_n^{-1}\}$ . Project back down to the real axis with  $\{z_n, z_n^{-1}\} \mapsto \frac{1}{2}(z_n + z_n^{-1}) = x_n$ .
- Toss out any “roots”  $x_n \notin [-1, 1]$ .

# *Demo*

# Future Work

The real power of Chebyshev functions is in creating differential operators and solving differential equations: given  $f \in L^2[-1, 1]$  solve

$$Lu = f.$$

- Reduces to a dense  $(N \times N)$  matrix solve.
- Distinction from finite difference methods: spectral accuracy is preserved! This means  $O(\epsilon_{\text{mach}})$  accurate solutions,  $u$ .
- You can implement this! (See next slide.)



## Future Work

The obvious issue: don't want to compete with Trefethen et. al.'s work. (Unless I develop an awesomer algorithm!) Possible solutions:

- 1 parallel projects in Python/Sage and Matlab,
- 2 a core C/FORTRAN library with separate Python and Matlab interfaces. (A la LAPACK.)

Current Python implementation is an excellent platform for new Python programmers. Learn Python and Mercurial! Get involved!

<http://code.google.com/p/pychebfun>

# Thank You